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# Some Applications of Spanning Trees in Complete and Complete Bipartite Graphs

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**Abstract: Problem statement:** The number of spanning trees  $\tau(G)$  in graphs (networks) is an important invariant, it is also an important measure of reliability of a network. **Approach:** Using linear algebra and matrix analysis techniques to evaluate the associated determinants. **Results:** In this study we derive simple formulas for the number of spanning trees of complete graph  $K_n$  and complete bipartite graph  $K_{n,m}$  and some of their applications. A large number of theorems of number of the spanning trees of known operations on complete graph  $K_n$  and complete bipartite graph  $K_{n,m}$  are obtained. **Conclusion:** The evaluation of number of spanning trees is not only interesting from a mathematical (computational) perspective, but also, it is an important measure of reliability of a network and designing electrical circuits. Some computationally hard problems such as the travelling salesman problem can be solved approximately by using spanning trees. Due to the high dependence of the network design and reliability on the graph theory we introduced the following important theorems and lemmas and their proofs.

Key words: Complete graph, complete bipartite graph, spanning trees, Kirchhoff matrix, operations on graphs

### INTRODUCTION

We consider finite undirected graph with no loops or multiple edges. Let G be such a graph on n vertices. A spanning tree for a graph G is a subgraph of G that is a tree and contains all vertices of G. There are many situation in which good spanning trees must be found. Whenever one wants to find a simple, cheap, yet efficient way to connect a set of terminals, be they computers, telephones, factories, or cities, a solution is normally one kind of spanning trees. Spanning trees prove important for several reasons: They create a spare subgraph that reflects a lot about the original graph, they play an important role to designing efficient routing algorithms, some computationally hard problems, such as the Steiner tree problem and the travelling salesperson problem, can be solved by using spanning trees and they have wide applications in many areas such as network design, bioinformatics (Biggs, 1993; Brown et al., 1996; Colbourn, 1987; Bermond et al., 1995; Myrvold et al., 1991). The number of spanning trees of G, denoted by  $\tau(G)$ , is the total number of distinct spanning subgraphs of G that are trees. A classic result of Kirchhoff, (Cayley, 1889) can be used to determine the number of spanning trees for G = (V, E). Let  $V = v_1, \dots, v_n$ . To state the result, we define the  $n \times n$  characteristic matrix  $A = [a_{ij}]$  as follows: (i)  $a_{ij} = -1$  if  $v_i$  and  $v_j$  are adjacent and  $i \neq j$ , (ii)  $a_{ii}$  equals the degree of vertex  $v_i$  if i = j and (iii)  $a_{ii} = 0$ otherwise. The Kirchhoff matrix tree theorem states that all cofactors of A are equal and their common value is  $\tau(G)$ . The matrix tree theorem can be applied to any graph G to determine  $\tau(G)$ , but this requires evaluating a determinant of a corresponding characteristic matrix. However, for a few special families of graphs there exists simple formulas that make it much easier to calculate and determine the number of corresponding spanning trees especially when these numbers are very large. One of the first such result is due to Cayley who showed that complete graph on n vertices, K<sub>n</sub> has n<sup>n-2</sup> spanning trees (Cvetkovic et al., 1980) that he showed  $\tau(K_n) = n^{n-2}$ ,  $n \ge 2$ . Another result,  $\tau(K_n) = p^{q-1}$ ,  $q^{p-1}p,q \ge 1$ , where  $K_{p,q}$  is the complete bipartite graph with bipartite sets containing p and q vertices,

Corresponding Author: Daoud, S.N., Department of Mathematics, Faculty of Science, El-Minufiya University, Shebeen El-Kom, Egypt respectively. It is well known, as in e.g., (Petingi *et al.*, 1998; Austin, 1960; Clark, 2003; Egecioglu and Remmel, 1994; Porter, 2004; Lewis, 1999).

Let  $G_1$  and  $G_2$  be a simple graphs. We introduce some operations on graphs:

- If G<sub>1</sub> and G<sub>2</sub> are vertex disjoint graphs. Then the join product, G<sub>1</sub>∨G<sub>2</sub>, is the super-graph of G<sub>1</sub> + G<sub>2</sub>, in which each vertex of G<sub>1</sub> is adjacent to every vertex of G<sub>2</sub>, (Balakrishnan and Ranganathan, 2000).
- The Cartesian product of two graphs  $G_1$  and  $G_2$ ,  $G_{1\times} G_2$ , is the simple graph with vertex set  $V(G_{1\times}$   $G_2) = V (V_1 \times V_2)$  and edge set  $E(G_1 \times G_2) = [(E_1 \times V_2) \cup (V_1 \times E_1)]$ , such that two vertices  $(u_1, u_2)$  and  $(v_1, v_2)$  are adjacent in  $G_1 \times G_2$  iff, either  $u_1 = v_1$  and  $u_2$  is adjacent to  $v_2$  in  $G_1$ , or  $u_1$  is adjacent to  $v_1$  in  $G_1$  and  $u_2 = v_2$ , (Wilson and Watkins, 1990).
- The tensor product, or Kronecher product of two graphs G₁ and G₂, G₁ ⊗ G₂, is the simple graph with v (G₁ ⊗ G₂) = V₁×V₂ where (u₁, u₂) and (v₁, v₂) are adjacent in G₁ ⊗ G₂ iff u₁ is adjacent to v₁ in G₁ and u₂ is adjacent to v₂ in G₂, (Balakrishnan and Ranganathan, 2000).
- The normal product, or the strong produc of two graphs  $G_1$  and  $G_2$  t,  $G_1 \circ G_2$ , is the simple graph with  $V(G_1 \circ G_2) = V_1 \times V_2$  where  $(u_1, u_2)$  and  $(v_1, v_2)$  are adjacent in  $G_1 \circ G_2$  iff either  $u_1 = v_1$  and  $u_2$  is adjacent to  $v_2$  in  $G_2$ , or  $u_1$  is adjacent to  $v_1$  in  $G_1$  and  $u_2 = v_2$ , or  $u_1$  is adjacent to  $v_1$  and  $u_2$  is adjacent to  $v_2$  (Balakrishnan and Ranganathan, 2000).
- The corona G<sub>1</sub> ⊙G<sub>2</sub> of G<sub>1</sub> and G<sub>2</sub> is the graph obtained by taking one copy of G<sub>1</sub> (which has n<sub>1</sub> vertices) and n<sub>1</sub> copies of G<sub>2</sub> and then joining the i<sup>th</sup> vertex of G<sub>2</sub> to every vertex in the i<sup>th</sup> copy of G<sub>2</sub>, (Wilson and Watkins, 1990).

The well-known matrix tree theorem (Kirchhoff matrix) can be used to count the number of spanning trees for small graphs, but this method is not feasible for large graphs. So we present two formulas in lemma1, lemma2 is to express  $\tau(G)$  directly as a determinant rather than in terms of cofactors as in Kirchhoff theorem.

**Lemma 1:** Let G be a simple graph with n vertices. Then,  $\tau(G) = \frac{1}{n^2} \det[D - A + U]$ , where A,D are the adjacency and degree matrices of G respectively and U is the n×n matrix where all its elements are ones. **Proof:** By simple calculations using the property of addition of two determinants, we can write det[D-A+U] as an addition of n determinants each of which is the same as det[D-A], but by replace one of its column and any determinant of these is equal to n× det[D-A]. Since all cofactors of [D-A] are equals. Then we have:

$$det[D - A + U] = n \times cofactor[D - A] +$$
  
n \times cofactor[D - A] + ..... + n \times cofactor  
[D - A](n - times) + det[D - A]

But:

$$det[D-A] = 0$$

Then:

$$det[D - A + U] = n^{2} \times cofactor[D - A] = n^{2}\tau(G)$$

Therefore:

$$\tau(G) = \frac{1}{n^2} \det[D - A + U]$$

**Lemma 2:** Let G be a simple graph with vertices n. Then,  $\tau(G) = \frac{1}{n^2} \det[nI_n - \overline{D} + \overline{A}]$ , where  $\overline{A}, \overline{D}$  are the adjacency and degree matrices of  $\overline{G}$  ( complement of G) respectively and  $I_n$  is the identity matrix.

**Proof:** It is clear that  $D + \overline{D} = (n-1)I_n$  and  $(U - A - I_n) = \overline{A}$ . Then.  $D - A + U = (n-1)I_n - \overline{D} + \overline{A} + I_n = nI_n - \overline{D} + \overline{A}$ . Thus from lemma1, we have:  $\tau(G) = \frac{1}{n^2} det[nI_n - \overline{D} + \overline{A}]$ .

**Theorem 3:**  $\tau(K_n) = n^{n-2}$ .

**Proof:** Applying lemma 2, we have:

$$\begin{aligned} \tau(K_n) &= \frac{1}{n^2} \det(n I - \overline{D} + \overline{A}) = \\ & \frac{1}{n^2} \det\begin{pmatrix} n & 0 & 0 & \cdots & 0 \\ 0 & n & 0 & \ddots & \vdots \\ 0 & \ddots & \ddots & \ddots & 0 \\ \vdots & \ddots & \ddots & n & 0 \\ 0 & \cdots & 0 & 0 & n \end{pmatrix} = \frac{1}{n^2} \times n^n = n^{n-2} \end{aligned}$$

**Corollary 4:**  $\tau(K_n - e) = (n - 2)n^{n-3}$ .

**Proof:** Applying lemma 2, we have:

$$\begin{aligned} \tau(K_n - e) &= \frac{1}{n^2} \det(n I - \overline{D} + \overline{A}) \\ &= \frac{1}{n^2} \det\begin{pmatrix} n - 1 & 1 & 0 & \cdots & 0\\ 1 & n - 1 & 0 & \ddots & \vdots\\ 0 & \ddots & \ddots & \ddots & 0\\ \vdots & \ddots & \ddots & n & 0\\ 0 & \cdots & 0 & 0 & n \end{pmatrix} \\ &= \frac{1}{n^2} \times ((n - 1)^2 - 1) \times n^{n-2} = (n - 2) \times n^{n-3} \end{aligned}$$

**Corollary 5:**  $\tau(K_n \circ e) = 2n^{n-3}$ , where  $K_n \circ e$  is the graph obtained from  $K_n$  by contracting the edge e.

**Proof:** Immediately from the fact that:

$$\tau(G) = \tau(G - e) + \tau(G \circ e)$$

**Theorem 6:** Let G be a graph constructed by removing m distinct edges from  $K_n, n \ge 2m$ . Then:

$$\tau(\mathbf{G}) = n^{n-2} (1 - \frac{2}{n})^m$$

**Proof:** Straightforward induction using properties of determinants.

**Theorem 7:**  $\tau(K_n + e) = (n+2)n^{n-3}$ .

**Proof:** Applying lemma 2, we have:

$$\tau(K_n + e) = \frac{1}{n^2} \det(nI - \overline{D} + \overline{A}) =$$

$$\frac{1}{n^2} \det\begin{pmatrix} n+1 & -1 & 0 & \cdots & 0\\ -1 & n+1 & 0 & \ddots & \vdots\\ 0 & \ddots & n & \ddots & 0\\ \vdots & \ddots & \ddots & \ddots & 0\\ 0 & \cdots & 0 & 0 & n \end{pmatrix}$$

$$= \frac{1}{n^2} \times ((n+1)^2 - 1) \times n^{n-2} = (n+2) \times n^{n-3}$$

**Theorem 8:** Let  $H_n$  be a graph constructed by removing n distinct edges from  $K_{2n}$ . Then:

$$\tau(H_n) = 2^{2n-2} \times n^{n-1} \times (n-1)^n$$

**Proof:** Applying lemma 2, we have:

$$\tau(H_n) = \frac{1}{(2n)^2} \det(2nI - \overline{D} + \overline{A}) = \frac{1}{(2n)^2} \det$$

2n - 1	1	0				•••	•••	•••	0
1	2n-1	0	۰.	÷	·	۰.	۰.		÷
0	0	2n - 1	1	0	÷	·	۰.	·	÷
÷	·	1	2n - 1	0	·			٠.	÷
÷		0	0	2n - 1	1	0			÷
÷	÷		۰.	1	2n - 1	0			÷
÷	÷	·		·	0	·	0	٠.	÷
	·	·	۰.	÷	0	1	۰.		0
		·	۰.	·	÷	·	۰.	2n-1	1
0							0	1	2n-1

Straightforward induction using properties of determinants. We have:

$$\tau(H_n) = \frac{1}{(2n)^2} \times ((2n-1)^2 - 1)^n = 2^{2n-2} \times n^{n-1} \times (n-1)^n$$

**Theorem 9:** Let G be a graph constructed by removing a star graph  $K_{1,2}$  from  $K_n$ . Then:

$$\tau(G) = n^{n-2}(1-\frac{1}{n})(1-\frac{3}{n})$$

**Proof:** Apply lemma 2, we have:

$$\begin{split} \tau(G) &= \frac{1}{n^2} \det(n \ I - \overline{D} + \overline{A}) = \frac{1}{n^2} \det\\ \begin{pmatrix} n-2 & 1 & 1 & 0 & \cdots & \cdots & \cdots \\ 1 & n-1 & 0 & \ddots & \vdots & \ddots & \ddots \\ 1 & 0 & n-1 & 0 & 0 & \vdots & \ddots & \cdot \\ 0 & \ddots & 0 & n & 0 & \ddots & \cdots & \\ \vdots & \cdots & 0 & 0 & \ddots & 0 & 0 \\ \vdots & \vdots & \ddots & \ddots & 0 & n \\ \end{pmatrix} \\ &= \frac{1}{n^2} \det \begin{pmatrix} n-2 & 1 & 1 \\ 1 & n-1 & 0 \\ 1 & 0 & n-1 \end{pmatrix} \times \det \begin{pmatrix} n & 0 & \cdots \\ 0 & n & \ddots \\ \vdots & \ddots & n \end{pmatrix} \\ &= \frac{1}{n^2} \times n^{n-3} \times (n-1)(n^2 - 3n) = n^{n-2} \times (1 - \frac{1}{n}) \times (1 - \frac{3}{n}) \end{split}$$

**Corollary 10:** Let G be a graph constructed by removing a star graph  $K_{l,m}$  from Kn,  $n \ge m+1$ .. Then:

$$\tau(G) = n^{n-2} (1 - \frac{1}{n})^{m-1} (1 - \frac{m+1}{n})$$

**Proof:** Straightforward induction using properties of determinants.

**Lemma 11:** Biggs (1993) let G be a k-regular graph with n vertices and m edges. Then:

$$\tau(L(G)) = 2^{m-n+1} \times k^{m-n-1} \times \tau(G)$$

where, L(G) is the line graph of G. **Theorem 12:** Let  $T_n$  be the line graph of  $K_n$ . Then:

$$\tau(T_n) = \tau(L(K_n)) =$$
  
$$2^{\frac{1}{2}(n^2 - 3n + 2)} \times (n - 1)^{\frac{1}{2}(n^2 - 3n - 2)} \times n^{n-2}$$

**Proof:** The line graph L(G) of a graph G is constructed by taking the edges of G as vertices of L(G) and joining two vertices in L(G) whenever the corresponding edges in G have a common vertex. Also if G is regular of valency k, its line graph L(G) is regular of valency 2k-2. It is easy to show that L(K<sub>n</sub>) is the triangle graph T<sub>n</sub> which can be described by saying that the  $\frac{1}{2}n(n-1)$  pairs of numbers from the set {1,2,...,n}, two vertices being adjacent whenever the corresponding pairs have just one common member. Appling lemma11 taking k = n-1,  $m = \frac{1}{2}n(n-1)$ , we have:

$$\tau(T_n) = \tau(L(K_n)) = 2^{\frac{1}{2}(n^2 - 3n + 2)} \times (n - 1)^{\frac{1}{2}(n^2 - 3n - 2)} \times n^{n - 2}$$

**Lemma 13:** Let  $A_n(x)$  be  $n \times n$  matrix such that:

$$A_{n}(x) = \begin{pmatrix} x & 1 & 1 & \cdots & \cdots & 1 \\ 1 & x & 1 & \ddots & & \vdots \\ 1 & \ddots & \ddots & 1 & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 \\ \vdots & & \ddots & \ddots & x & 1 \\ 1 & \cdots & \cdots & 1 & 1 & x \end{pmatrix}$$

Then:

$$\det(A_{n}) = (x + n - 1)(x - 1)^{n}$$

**Proof:** From the definition of the circulant determinants, we have:

$$det(A_n(x)) = \prod_{j=1}^n (x + \omega_j + \omega_j^2 + \omega_j^3 + \dots + \omega_j^{n-1})$$
$$= (x + 1 + 1 + \dots + 1) \times \prod_{j=1,\omega_j \neq 1}^n$$
$$(x + \underbrace{\omega_j + \omega_j^2 + \omega_j^3 + \dots + \omega_j^{n-1}}_{=-1})$$
$$= (x + n - 1) \times (x - 1)^{n-1}.$$

**Lemma14:** If 
$$H = \begin{pmatrix} A & B \\ B & A \end{pmatrix}$$
 and  $AB = BA$ . Then  $det(H) = det(A+B)$ .  $det(A-B)$ .

Proof: Using the fact that

$$det \begin{pmatrix} A & B \\ C & D \end{pmatrix} = \begin{cases} det(A).det(D - CA^{-1}B), \text{ where } A, B \\ det(D).det(A - BD^{-1}C) \end{cases}$$
  
are non singular, Marcus M. [12]. We have  
$$det \begin{pmatrix} A & B \\ B & A \end{pmatrix} = det(A).det(A - BA^{-1}B)$$
  
$$= det(A^2 - B^2) = det(A + B).det(A - B).$$

This formula gives some sort of symmetry in some matrices which facilitates our calculation of determinants.

**Theorem 15:**  $\tau(K_2 \times K_n) = n^{n-2} \times (n+2)^{n-1}$ . **Proof:** Applying lemmas 2. We have:

$\tau(K_2)$	$\langle K_n \rangle =$	$=\frac{1}{(2r)}$	$\frac{1}{n^2}$ det(	(2n I –	$\overline{D} + \overline{A}$	$=\frac{1}{(2n)}$	$\frac{1}{(1-1)^2}$ de	et	
(n+1	0			0	0	1			1
0	n + 1	0	·	:	1	0	1		:
:	0	·	·	:	:	·	·.	·.	:
:	·.	·.	n + 1	0	:		1	0	1
0			0	n + 1	1	•••		1	0
0	1			1	n + 1	0			0
1	0	1		÷	0	n + 1	0	·	÷
	·.	·.	·	÷	:	0	·		÷
		·.	0	1	÷	·	·	n + 1	0
1			1	0	0			0	n+1

Using lemma14, we get:

$$\mathbf{t}(\mathbf{K}_{2} \times \mathbf{K}_{n}) = \frac{1}{(2n)^{2}} \det \begin{pmatrix} n+1 & 1 & 1 & \cdots & 1\\ 1 & n+1 & 1 & \ddots & \vdots\\ 1 & 1 & \ddots & \ddots & 1\\ \vdots & \ddots & \ddots & n+1 & 1\\ 1 & \cdots & 1 & 1 & n+1 \end{pmatrix} \times \det \begin{pmatrix} n+1 & -1 & \cdots & \cdots & -1\\ -1 & n+1 & -1 & \ddots & \vdots\\ \vdots & \ddots & \ddots & \ddots & \vdots\\ \vdots & \ddots & \ddots & n+1 & -1\\ -1 & \cdots & \cdots & -1 & n+1 \end{pmatrix}$$

Applying lemma13 with x = n + 1, for the first determinant and properties of determinants for the second, yields:

$$\tau(K_2 \times K_n) = \frac{1}{(2n)^2} \times 2n^n \times 2(n+2)^{n-1} = (n+2)^{n-1} \times n^{n-2}$$

**Theorem 16:**  $\tau(K_2 \otimes K_n) = n^{n-2}(n-1)(n-2)^{n-1}$ .

**Proof:** Applying lemmas 2. We have:

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$$\tau(\mathbf{K}_{2} \otimes \mathbf{K}_{n}) = \frac{1}{(2n)^{2}} \det(2n\mathbf{I} - \overline{\mathbf{D}} + \overline{\mathbf{A}}) =$$

$$\frac{1}{(2n)^{2}} \det\begin{pmatrix}n & 1 & \cdots & 1\\ 1 & \ddots & \vdots & & \mathbf{I}_{n}\\ & \ddots & 1 & & & \\ 1 & \cdots & 1 & n & & \\ & & & 1 & \ddots & & \\ & & & & 1 & \ddots & & \\ & & & & 1 & \cdots & 1 & n \end{pmatrix}$$

Using lemma14, we get:

$$\begin{split} \tau(\mathbf{K}_2\otimes\mathbf{K}_n) &= \\ \frac{1}{(2n)^2} \det \begin{pmatrix} n+1 & 1 & 1 & \cdots & 1\\ 1 & n+1 & 1 & \ddots & \vdots\\ 1 & 1 & \ddots & \ddots & 1\\ \vdots & \ddots & \ddots & n+1 & 1\\ 1 & \cdots & 1 & 1 & n+1 \end{pmatrix} \times \det \begin{pmatrix} n-1 & 1 & 1 & \cdots & 1\\ 1 & n-1 & 1 & \ddots & \vdots\\ 1 & 1 & \ddots & \ddots & 1\\ \vdots & \ddots & \ddots & n-1 & 1\\ 1 & \cdots & 1 & 1 & n-1 \end{pmatrix} \end{split}$$

Applying lemma13 with x = n + 1, n - 1 respectively, we have:

$$\tau(K_2 \otimes K_n) = \frac{1}{(2n)^2} \times 2n^n \times 2(n-1)(n-2)^{n-1}$$
$$= n^{n-2}(n-1)(n-2)^{n-1}$$

**Theorem 17:**  $\tau(K_2 \circ K_n) = (2n)^{2n-2}$ .

## **Proof:** Applying lemma 2, we have:

$$\begin{split} \tau(\mathbf{K}_{2} \circ \mathbf{K}_{n}) &= \frac{1}{(2n)^{2}} \det(2n \, \mathbf{I} - \overline{\mathbf{D}} + \overline{\mathbf{A}}) = \frac{1}{(2n)^{2}} \det(2n \, \mathbf{I} - \mathbf{D} + \overline{\mathbf{A}}) \\ & \begin{pmatrix} 2n & 0 & \cdots & \cdots & 0 & 0 & \cdots & \cdots & \cdots & 0 \\ 0 & 2n & 0 & \ddots & \vdots & \vdots & 0 & \cdots & \cdots & \vdots \\ \vdots & 0 & \ddots & \ddots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \\ 0 & \cdots & \cdots & 0 & 2n & 0 & \cdots & \cdots & 0 & 0 \\ 0 & \cdots & \cdots & 0 & 2n & 0 & \cdots & \cdots & 0 \\ \vdots & 0 & \cdots & \cdots & 0 & 2n & 0 & \cdots & \cdots & 0 \\ \vdots & 0 & \cdots & \cdots & \vdots & 0 & 2n & 0 & \ddots & \vdots \\ & \ddots & \ddots & \vdots & \vdots & 0 & \ddots & \vdots \\ & \ddots & \ddots & \vdots & \vdots & 0 & \ddots & \vdots \\ & 0 & \cdots & \cdots & 0 & 0 & \cdots & \cdots & 0 & 2n \\ \end{pmatrix} \\ & & = \frac{1}{(2n)^{2}} (2n)^{n} \times (2n)^{n} = (2n)^{2n-2} = \tau(\mathbf{K}_{2n}) \end{split}$$

**Theorem 18:**  $\tau(K_2 \odot K_n) = n^{2(n-2)}$ .

**Proof:** Applying lemma 2, we have:

$$\tau(K_2 \odot K_n) = \frac{1}{(2n)^2} \det(2nI - \overline{D} + \overline{A}) = \frac{1}{(2n)^2} \det$$

n + 1	0	•••	•••	0	0	1	•••	•••	1)
0	n	0	·.	÷	1	1			:
÷	0	·.	·.	÷	÷	·.	·.		:
÷	·.	·.	n	0	÷		·.	·	:
0			0	n	1				1
0				1	m 1 1	Ο			0
0	1	•••	•••	1	$\Pi \pm I$	0			0
0 1	1 1		····	:	11 + 1	n	0	·	:
0 1 :	1 1 ·.	···· ··· ··	···· ··· ·.	:	$ \begin{array}{c} 0\\ \vdots \end{array} $	n 0	0 ·.	·	:
0 1 : :	1 1 ·.	···· ··· ··. ··.	···· ··. ··.	1 : :	0 : :	n 0 ∵.	0 · ·	∴. n	0 : : 0
0 1 : : 1	1 1 ·	···· ··. ··. ··.	···· ··. ···	1 : : : 1	0 : : 0	n 0 ·	0 · ·.	· n 0	0 : : 0 n

Using lemma14, we get:

$$\tau(\mathbf{K}_{2} \odot \mathbf{K}_{n}) = \frac{1}{(2n)^{2}} \det \begin{pmatrix} n+1 & 1 & 1 & 1 \\ 1 & n+1 & 1 & \vdots \\ 1 & 1 & \ddots & \vdots \\ \vdots & \cdots & \ddots & n+1 \end{pmatrix} \times \det \begin{pmatrix} n+1 & -1 & \cdots & -1 \\ -1 & n-1 & -1 & \ddots \\ \vdots & \ddots & \ddots & -1 \\ -1 & \cdots & -1 & n-1 \end{pmatrix}$$

Applying lemma13 with x = n+1 for the first determinant and properties of determinants for the second, yields:

$$\tau(K_2 \odot K_n) = \frac{1}{(2n)^2} \times 2n^n \times 2n^{n-2} = n^{2(n-2)}$$

**Theorem 19:**  $\tau(K_{n,m}) = n^{m-1}m^{n-1}$ .

**Proof:** Apply lemma 2, we have:

$$\tau(\mathbf{K}_{n,m}) = \frac{1}{(n+m)^2} \det((n+m)\mathbf{I} - \overline{\mathbf{D}} + \overline{\mathbf{A}}) = \frac{1}{(n+m)^2} \det((n+m)\mathbf{I} - \overline{\mathbf{D}} + \overline{\mathbf{A}}) = \frac{1}{(n+m)^2} \det((n+m)^2) \det((n+m)^2)$$

Using lemma14, we get:

$$\tau(\mathbf{K}_{n,m}) = \frac{1}{(n+m)^2} \det \begin{pmatrix} m+1 & 1 & \cdots & 1 \\ 1 & \ddots & \vdots \\ \vdots & \ddots & 1 \\ 1 & \cdots & 1 & m+1 \end{pmatrix} \times \det \begin{pmatrix} n+1 & 1 & \cdots & 1 \\ 1 & \ddots & \vdots \\ \vdots & \ddots & 1 \\ 1 & \cdots & 1 & n+1 \end{pmatrix}$$

Applying lemma13 with x = n + 1 and x = m + 1 respectively, we have:

$$\tau(K_{n,m}) = \frac{1}{(n+m)^2} \cdot (n+m) \times m^{n-1} \times (n+m) \times n^{m-1} = m^{n-1} n^{m-1}$$

Specially,  $\tau(K_{n,n}) = n^{2n-2}$ .

**Corollary 20:**  $\tau(K_{n,n} - e) = n^{2n-4} \times (n-1)^2$ .

**Proof:** Applying lemma 2, we have:

τ(Κ	( <sub>n,n</sub> – e)	$=\frac{1}{(2)}$	$\frac{1}{n^2}$ det(	2n I	$-\overline{D}+\overline{A}$	$\bar{A}$ ) = -	$\frac{1}{(2n)^2}$ det
(n	1		1	1	0		0)
:	n + 1	÷	÷	0	0		:
:	۰.	۰.	1	÷			:
1	٠.	1	n + 1	0			0
1	0		0	n	1		1
0	0		÷	1	n + 1	÷	:
:	•••		÷	÷	۰.	·.	1
0			0	1		1	n + 1)

Using lemma14, we get:

$$\tau(K_{n,n} - e) = \frac{1}{(2n)^2} \det \begin{pmatrix} n+1 & 1 & \cdots & 1 \\ 1 & \ddots & \vdots \\ \vdots & \ddots & 1 \\ 1 & \cdots & 1 & n+1 \end{pmatrix} \times \det \begin{pmatrix} n-1 & 1 & \cdots & 1 \\ 1 & n+1 & \vdots \\ \vdots & & \ddots & 1 \\ 1 & \cdots & 1 & n+1 \end{pmatrix}$$

Applying lemma13 with x = n+1 for the first determinant and properties of determinants for the second, yields:

$$\tau(\mathbf{K}_{n,n} - \mathbf{e}) = \frac{1}{(2n)^2} \times 2n^n \times 2n^{n-2} \times (n-1)^2 = n^{2n-4} \times (n-1)^2$$

**Theorem 21:**  $\tau(K_{n,n} + e) = n^{2(n-2)}(n^2 + 2n - 1)$ .

**Proof:** Applying lemma 2, we have:

$$\tau(K_{n,n} + e) = \frac{1}{(2n)^2} \det(2nI - \overline{D} + \overline{A}) = \frac{1}{(2n)^2} \det(2nI - \overline{D} + \overline{A})$$

(n+2)	1			1	$^{-1}$	0			0)
1	n + 1	1	·	÷	0	0			:
:	1	·	·	:	:	·	·	·	:
:	·.	·	n+1	1	:		·	·	:
1			1	n + 1	0				0
-1	0			0	n+2	1			1
0	0			÷	1	n + 1	1	·	:
1 .									
:	· .	·		÷	-	1	·•.		:
:	·	•	·. 	:	:	1 •.	•	n + 1	1

Using lemma14, we get:

$$\begin{split} \tau(\mathbf{K}_{n,n}+e) = \\ \frac{1}{(2n)^2} det \begin{pmatrix} n+1 & 1 & 1 & \cdots & 1\\ 1 & n+1 & 1 & \ddots & \vdots\\ 1 & 1 & \ddots & \ddots & 1\\ \vdots & \ddots & \ddots & n+1 & 1\\ 1 & \cdots & 1 & 1 & n+1 \end{pmatrix} \times det \begin{pmatrix} n+3 & 1 & 1 & \cdots & 1\\ 1 & n+1 & 1 & \ddots & \vdots\\ 1 & 1 & \ddots & \ddots & 1\\ \vdots & \ddots & \ddots & n+1 & 1\\ 1 & \cdots & 1 & 1 & n+1 \end{pmatrix} \end{split}$$

Applying lemma13 with x = n+1 for the first determinant and properties of determinants for the for the second, yields:

$$\tau(\mathbf{K}_{n,n} + \mathbf{e}) = \frac{1}{(2n)^2} \times 2n^n \times 2n^{n-2}(n^2 + 2n - 1)$$
$$= n^{2(n-2)}(n^2 + 2n - 1)$$

**Theorem 22:**  $\tau(L(K_{n,n})) = 2^{n^2-2n+1} \times n^{n^2-2n-1} \times n^{2(n-1)}$ .

**Proof:** It is easy to show that  $L(K_{n,n})$  is the  $K_n \times K_n$ .

Appling lemma11 taking k = n and  $m = n^2$ . We have:

$$\tau(L(K_{n,n})) = \tau(K_n \times K_n) = 2^{n^2 - 2n + 1} \times n^{n^2 - 2n - 1} \times n^{2(n-1)}$$

Theorem23:

 $\tau(K_1 + K_{n,m}) = (n + m + 1) \times (n + 1)^{m-1} \times (m + 1)^{n-1}$ .

**Proof:** Applying lemma 2, we have:

$$\tau(\mathbf{K}_{1} + \mathbf{K}_{n,m}) = \frac{1}{(n+m+1)^{2}} \det[(n+m+1)\mathbf{I} - \overline{\mathbf{D}} + \overline{\mathbf{A}}]$$

$$= (n+m+1) \times \frac{1}{(n+m+1)^{2}} \det$$

$$\begin{pmatrix} m+2 & 1 & \cdots & 1 & 0 & \cdots & \cdots & 0 & 0 \\ \vdots & \ddots & \vdots & \vdots & \vdots & \cdots & \cdots & \vdots & \vdots \\ 1 & \ddots & 1 & m+2 & 0 & \cdots & \cdots & 0 & \vdots \\ 0 & \cdots & \cdots & 0 & n+2 & 1 & \cdots & 1 & \vdots \\ \vdots & \cdots & \cdots & \vdots & 1 & \ddots & \vdots & \vdots & \vdots \\ \vdots & \cdots & \cdots & \vdots & 1 & \ddots & \vdots & \vdots & \vdots \\ \vdots & \cdots & \cdots & \vdots & \vdots & \ddots & \ddots & 1 & \vdots \\ 0 & \cdots & \cdots & 0 & 1 & \cdots & 1 & n+2 & 0 \\ 0 & \cdots & \cdots & \cdots & \cdots & \cdots & \cdots & n+m+1 \end{pmatrix}$$

= (n	$= (n+m+1) \times \frac{1}{(n+m+1)^2} \det$												
( m -	+ 2	1		1		(n+2	1		1)				
1		٠.		÷		1	n + 2		:				
:			۰.	1	× det	:		۰.	1				
			1	m + 2)	)	1		1	n+2				

Applying lemma 13 with x = m + 2 and x = n + 2, respectively, we have:

 $\tau(K_{1} + K_{n,m}) = (n + m + 1) \times (n + 1)^{m-1} \times (m + 1)^{n-1}$ 

Specially,  $\tau(K_1 + K_{n,n}) = (2n+1) \times (n+1)^{2n-2}$ .

**Corollary24:**  $\tau(K_{n,n} \circ e) = (2n-1)n^{2n-4} = \tau(K_1 + K_{n-1,n-1})$ .

**Theorem 25:**  $\tau(K_2 \times K_{n,n}) = 2n^{2n-2} \times (n+1) \times (n+2)^{2n-2}$ .

**Proof:** Applying lemma 2, we have:

1	(K	$_2 \times K$	.,,n)	=-(	1 4n	$\frac{1}{\left(1-\frac{1}{2}\right)^2}$ de	t(4n	I – D	) + Ā	)=-(	$\frac{1}{4n)^2}d$	let	
(	n+2	0	1	0	1			0	1	1			1
ł	0	n+2	0	1	0	·.	·	1	0	1	··. ··.	·	÷
	1	0	n+2	0	1	0	·.	÷	1	·	··. ··.	۰.	÷
ł	0	1	0	۰.	0	·.	·	÷	÷	·	··. ··.	·	÷
	1		·	·	·.	1	1	÷	÷	·	··. ··.	۰.	÷
Ì	÷			1	0	n+2	0	÷	÷	·	··. ··.	0	1
ł	÷			÷	1	0	n+2	1				1	0
	0	1			•••		1	n+2	0	1	0 1		
ł	1	0	·	·	·	·.	÷	0	n+2	0	··. ··.	۰.	÷
	÷	1	·.	٠.	·.	۰.	÷	1	0	n+2	··. ··.	·	÷
Ì	÷	·	·	·	·.	·.	÷	0	1	0	··. ··.	۰.	÷
	÷	·	·	۰.	۰.	·.	÷	÷	·	·	··. ··.	·	1
1	÷	·.	·.	۰.	۰.	0	1	÷	·	·	··. ··.	n+2	0
	1	·.	·.	۰.	۰.	1	0	÷	·	·	··. ··.	0	n+2

## Using lemma14, we get:

$$\tau(K_{2} \times K_{n,n}) = \frac{1}{(4n)^{2}} det \begin{pmatrix} n+2 & 1 & 2 & 1 & 2 & \cdots & \cdots \\ 1 & n+2 & 1 & \ddots & \ddots & \ddots & \ddots \\ 2 & 1 & n+2 & \ddots & 2 & \ddots & \ddots \\ 1 & \ddots & \ddots & \ddots & \ddots & \ddots & \ddots \\ 2 & \ddots & \ddots & \ddots & \ddots & \ddots & \ddots \\ \vdots & \ddots & \ddots & \ddots & \ddots & n+2 & 1 \\ \vdots & \ddots & \ddots & \ddots & \ddots & n+2 & 1 \\ \vdots & \ddots & \ddots & \ddots & \ddots & n+2 & 1 \\ 0 & -1 & n+2 & \ddots & \ddots & \ddots & 1 & n+2 \end{pmatrix} \times det \begin{pmatrix} n+2 & -1 & 0 & -1 & \cdots & \cdots & \cdots \\ -1 & n+2 & -1 & \ddots & \ddots & \ddots & \ddots & 1 \\ 0 & -1 & n+2 & \ddots & \ddots & \ddots & \vdots \\ -1 & \ddots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \ddots & 0 \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \ddots & 0 \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & 0 \\ \vdots & \ddots & \ddots & \ddots & \ddots & -1 & n+2 & -1 \\ \vdots & \cdots & \cdots & \cdots & 0 & -1 & n+2 \end{pmatrix}$$

Straightforard induction using properties of determinants. We have:

$$\tau(\mathbf{K}_{2} \times \mathbf{K}_{n,n}) = \frac{1}{(4n)^{2}} \times 8n^{2n} \times 4(n+1)(n+2)^{2n-2}$$
$$= 2n^{2n-2} \times (n+1) \times (n+2)^{2n-2}$$

**Theorem 26**:  $\tau(K_2 \circ K_{n,n}) = 2^{4n-2} \times n^{2n-2} \times (n+1)^{2n}$ .

**Proof:** Apply lemma 2, we have:

	τ(K	$_{2} \circ \mathbf{K}$	n,n	) =	$=\frac{1}{(4n)}$	$\frac{1}{\left(1-\frac{1}{2}\right)^2}$ de	t(4n	I – D	+ Ā)	$=\frac{1}{4}$	$\frac{1}{(n)^2}$	let
2n+2	0	1	0	1			0	0	1	0 1		)
0	2n+2	0	1	0	۰.	۰.	0	0	0	1 0	1	·
1	0	2n+2	0	1	0	۰.	:	1	۰.	··. ··.	·	·
0	1	0	۰.	0	۰.	۰.	0	0	·.	··. ··.	·.	·.
1		۰.	۰.	٠.	1	1	1	÷	۰.	··. ··.	·	·
÷			1	0	2n+2	0	:	:	·	··. ··.	·	·.
÷			÷	1	0	2n+2	1	÷	۰.	··. ··.	·	·
0	0	1	0	1			2n+2	0	1	0 1		
0	0	0	1	0	·.		0	2n+2	0	··. ··.	·	·
1	0	0	0	۰.	۰.	۰.	1	0	2n+2	··. ··.	·	·.
0	1	·	۰.	·.	·	·	0	·	0	··. ··.	۰.	·.
1	0	۰.	۰.	۰.	·.	·.	·	۰.	·.	··. ··.	·	·.
0	1	۰.	۰.	۰.	·	·	·	·	·	·. ·.	2n+2	0
:	·	۰.	۰.	۰.	·	·	۰.	·	·	·. ·.	0	2n+2

Using lemma14, we get:

$$t(K_{2} \circ K_{n,n}) = \frac{1}{(4n)^{2}} det \begin{pmatrix} 2n+2 & 0 & 2 & 0 & 2 & \cdots & \cdots \\ 0 & 2n+2 & 0 & \ddots & \ddots & \ddots & \ddots \\ 2 & 0 & 2n+2 & \ddots & 2 & \ddots & \ddots \\ 0 & \ddots & \ddots & \ddots & \ddots & \ddots & \ddots \\ 2 & \ddots & \ddots & \ddots & \ddots & \ddots & \ddots \\ 2 & \ddots & \ddots & \ddots & \ddots & \ddots & \ddots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 2n+2 & 0 \\ \vdots & \ddots & \ddots & \ddots & \ddots & 0 & 2n+2 \end{pmatrix} \\ det \begin{pmatrix} 2n+2 & 0 & 0 & \cdots & \cdots & 0 \\ 0 & 2n+2 & 0 & \ddots & \ddots & \vdots \\ \vdots & 0 & 2n+2 & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 0 & \cdots & \cdots & \cdots & 0 & 2n+2 \end{pmatrix}$$

Straightforward induction using properties of determinants.

We have:  $\tau(K_2 \circ K_{n,n}) = 2^{4n-2} \times n^{2n-2} \times (n+1)^{2n}$ 

**Theorem 27:**  $\tau(K_2 \odot K_{n,n}) = (n+1)^{4n-4} \times (2n+1)^2$ .

**Proof:** Applying lemma 2, we have:

1	τ(K <sub>2</sub> )	⊙ K <sub>n,</sub>	$_{n}) = -\frac{1}{2}$	4(21	1 n +	$\frac{1}{1}^{2}$ de	et((4n	+2)I-	- D +	Ā) =	$\frac{1}{4(2n+1)}$	$(-1)^2$ d	et
(	2n+2	0	0				0	0	1	1			1)
	0	n+2	0	1	0	·.	÷	1	1	۰.	··. ··.	۰.	:
l	0	0	n+2	0	1	·.	÷	1	·.	·.	··. ··.	۰.	:
l	÷	1	0	۰.	0	·.	÷	:	۰.	۰.	··. ··.	۰.	:
	÷	0	·.	۰.	۰.	·.	1	:	۰.	۰.	··. ··.	۰.	:
l	÷	÷	۰.	۰.	0	n+2	0	÷	۰.	۰.	··. ··.	۰.	:
ł	0	÷	۰.	۰.	1	0	n+2	1				۰.	1
ļ	0	1	1		•••		1	2n+2	0	0	0		0
ł	1	1		۰.	۰.	۰.	÷	0	n+2	0	1 0		
	1	·	۰.	۰.	۰.	۰.	÷	0	0	n+2	0 .	۰.	·
l	÷	·	·.	۰.	۰.	·.	÷	:	1	0	··. ··.	۰.	·
ł	÷	·	·.	۰.	۰.	·.	÷	:	:	۰.	··. ··.	۰.	·
	:	·.	·.	۰.	۰.	·.	÷	÷	÷	·.	··. ··.	n+2	0
	1				•••		1	0	÷	·.	·. 1	0	n+2)

Using lemma14, we get:

$$\tau(K_{2} \odot K_{n,n}) = \frac{1}{4(2n+1)^{2}} \det \begin{pmatrix} 2n+2 & 1 & 1 & \cdots & \cdots & 1 \\ 1 & n+3 & 1 & 2 & 1 & \cdots & \cdots \\ 1 & 1 & n+3 & \ddots & \ddots & \ddots & \ddots \\ \vdots & 2 & \ddots & \ddots & \ddots & \ddots & \ddots \\ \vdots & 1 & \ddots & \ddots & \ddots & \ddots & \ddots \\ \vdots & \vdots & \ddots & \ddots & \ddots & n+3 & 1 \\ 1 & \vdots & \ddots & \ddots & \ddots & n+3 & 1 \\ 1 & \vdots & \ddots & \ddots & \ddots & n+3 & 1 \\ 1 & \vdots & \ddots & \ddots & \ddots & n+3 & 1 \\ 1 & \vdots & \ddots & \ddots & \ddots & n+3 & 1 \\ 1 & \vdots & \ddots & \ddots & \ddots & n+3 & 1 \\ 1 & \vdots & \ddots & \ddots & \ddots & n+3 & 1 \\ \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \ddots \\ \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \ddots \\ \vdots & \vdots & \ddots & \ddots & \ddots & \ddots & \ddots \\ \vdots & \vdots & \ddots & \ddots & \ddots & n+1 & -1 \\ -1 & \vdots & \cdots & \cdots & 0 & -1 & n+1 \end{pmatrix}$$

Straightforward induction using properties of determinants. We have:

$$\tau(\mathbf{K}_2 \odot \mathbf{K}_{n,n}) = \frac{1}{4(2n+1)^2} \times 2(n+1)^{2n-2} \times (2n+1)^{2n-2} \times (2n+1)^{2n-2} \times (2n+1) = (n+1)^{4n-4} \times (2n+1)^2$$

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