

# On the Computerization of African Languages

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**Abstract:** In this article, a computer tool for processing African languages has been designed. It is intended to be a contribution to the automatic processing of African languages. The current study is focused on West African languages where five main languages from Niger, two from Mali and one from Burkina Faso are considered. After a brief review of African languages processing, we designed a tool which uses minimum resources and operates essentially on a dictionary and the characteristics of the language alphabet. The dictionary is represented using a trie data structure. For the sake of application, the designed tool operates as a spell checker. To detect and correct spelling errors, the edit distance and the specificities of the language are used. Although they do not have processing tools, it was shown that existing tools for computerized languages can be adapted to African languages efficiently. To extend the designed tool to any African language, we only need to provide an appropriate dictionary and alphabet.

**Keywords:** Computerization, Natural Language Processing, Spell Checker, African Language, Computer Resource

## Introduction

The importance of Automatic natural language processing cannot be overemphasized. It has many industrial applications such as spell checking, parsing, text indexing and retrieval of information from the Internet, voice recognition and synthesis, vocal control of domestic robots, automated response systems and machine translation (Kukich, 1992; Pierre, 2006). Applications such as text editing are used by millions of people every day. They are present in all computer systems, internet search engines and electronic gadgets. To make their use more effective, processing tools like spell checking and grammar correction are integrated to these softwares and the main objective is to assist the user. While spell checking is concerned about detecting and correcting single word errors, parsing is concerned about grammatical errors detection and correction through a rigorous syntax analysis. Though syntax analyzers are more important, the first step towards the development of such programs is that of spell checkers where three different techniques are used which are non-word error detection, isolated-word error correction and context error detection and correction (Kukich, 1992). Such programs are generally designed for a given language. Though different processing techniques exist for other

languages like English and French, African languages do not yet have processing tools.

## Review of African Languages Processing

Africa is a continent with a very high linguistic diversity. Estimated at about 1500-2000 languages, four main groupings can be distinguished which are Afro-Asiatic, Nilo-Saharan, Niger-Saharan and Khoisan (see the AG Bell Association web site). All African languages are considered official languages of the African Union. Representing one third of the world's languages (Van Der and Gilles-Maurice, 2003), African languages are an important and irreplaceable component of the linguistic heritage of humanity and its ecolinguistic diversity. According to Osborn (2006), it is clear that African languages are not yet widely used in the content of computing applications or on the internet. He observed that African languages are represented on the web but not prominently as media of communication. A significant number of sites that treat African languages were given by Diki-Kidiri and Edema (2003) but with minimal content in the languages themselves.

The existence of resources is the first step in the computerization of a language (Chanard and Popescu-Belis, 2001). The majority of well-resourced languages have well-formed corpuses but this is not the case for African

languages. Despite the existence of different projects aimed at the computerization of African languages such as Pal and DILAF, the resources of these languages are still very scarce. Nevertheless, important resources like bilingual and editorial dictionaries exist and can be used in favour of African languages computerization. Although they have a lot of differences, we believe in the possibility of using the little resources they have in common to develop processing tools. From a vocal point of view, different tones can be found in the words of African languages. For example, Hausa words have high tones and low tones and one can observe a flexion of gender and number (Mijinguini and Naroua, 2012). Many African languages are currently used by major radio stations of the world such as BBC (UK), CRI (China), Deutsche Welle (Germany), IRIB (Iran), Radio Moscow (Russia), RFI (France) and VOA (USA). Unfortunately, the presence of African languages on Internet is very precarious even though they represent 30% of the languages of the world (Van Der and Gilles-Maurice, 2003). The current researches on African languages choose oral and written corpuses as a transitional alternative or build a corpus from the Web (Gilles-Maurice, 2002).

Another difficulty to overcome in the computerization of African languages is text entry. In fact, computer keyboards are designed for well-resourced languages and are not compatible with African languages. To enter texts in languages like Hausa, Fulfulde and Bambara requires special softwares. Unfortunately, this is the case for all African languages due to the presence of special characters. An evaluation of keyboard layouts for five languages from Niger (Fulfulde, Hausa, Kanuri, Songhai-Zarma, Tamasheq) recommended LLACAN which covers all the symbols of the alphabets of those languages, produces valid Unicode code and requires less buttons to press (Enguehard and Naroua, 2008).

The ability of processing a language with the famous word processing softwares like MS Word and OpenOffice Writer is another important issue. Various techniques have been developed for well-resourced languages for various purposes. However, all existing techniques are limited and inadequate in the case of African languages. Despite the scarcity of linguistic resources, it is possible to develop computer tools for African languages and improve them over the time with the possibility to create extensions for some popular softwares.

## Spell Checking Techniques

The main objective of spell checkers is to detect and correct errors. Their task is composed of three sub-tasks: Detecting errors, generating possible corrections and ranking suggested corrections. To achieve this, various techniques were invented. Each technique is related

either to non-word error correction, real-word error correction, or both. Spelling errors may be typographical, cognitive or phonetic. Typographical errors occur when the keys are pressed in the wrong order. Cognitive errors arise from ignorance of the correct spelling of the word. Phonetic errors are special cases of cognitive errors. A phonetic error refers to a wrong word that is pronounced the same way as the correct word. It was shown that in typed texts, 1 to 3% of the errors are spelling errors (Daniel and James, 2000). Damerau (1964) stated that 80% of these errors are related to insertion, deletion, substitution, or transposition.

Error detection is to find incorrect words in a text. A wrong word is then marked by the application in charge of spell checking. If the word is really wrong, an error is said to be detected. Many authors have made important contributions in this area like (Damerau 1964; Cyril, 1967; Peterson, 1980; Zamora et al., 1981; Laurent, 2001; Suzan, 2002; Pierre, 2006). The main techniques used for non-word error detection in a text are either based on analysis of n-grams, or dictionary lookup (Kukich, 1992). The techniques based on n-grams are to analyze each n-gram of a given input word and check its validity in a precompiled table. These techniques usually require a dictionary or corpus that's large enough to determine the statistics table of n-grams (Kukich, 1992). A dictionary is a collection of correct or acceptable words. The techniques based on the use of a dictionary or lexicon involve taking a word as input and verifying its existence in the dictionary. Any word that is not in the dictionary is then considered wrong (Kukich, 1992). A detection algorithm based on dictionary lookup is given by Peterson (1980). Some of the data structures used in spell checking are hash tables, binary search trees, tries and finite automata. One of the famous algorithms in this area is that of Aho and Corasick (1975). The algorithm is to move through an abstract data structure called dictionary that contains the words to search by reading the text characters one by one. The data structure is implemented efficiently, which ensures that each character of the text is read only once. Generally, the dictionary is represented using a trie. A trie may be seen as a representation of the transition function of a deterministic finite automaton. The algorithm has a linear complexity in the size of the text and search strings. Comparatively, techniques using n-grams derived from a dictionary provide less accuracy than those using all the information in the dictionary. But, the latter ones are time consuming depending on the data structure used to represent the dictionary. A comparative study showed that the hash table provides better performance than the AVL tree, the Red-Black tree and Skip list (Mark, 2009). A comparison of five data structures was performed for the Punjabi dictionary (Lehal and Singh, 2000). It concerned binary search tree,

trie, ternary search tree, multi-way tree and reduced memory method tree. As a result, the binary search tree was found to be the most suitable data structure in terms of memory usage and time. But it is limited when it comes to suggest a list of candidates for the correction or find all words that differ by one or two characters. This limitation may be avoided by the use of a trie which offers almost the same time complexity with a binary search tree. Hash table and trie are shown to be the most suitable data structures for dictionary representation.

Error correction refers to the fact of equipping spell checkers with the ability to correct detected errors. This is to find words in the dictionary that are similar in some ways to the misspelled word. The minimum edit distance or simply edit distance is until now the most widely used technique in the spelling errors correction. It has been applied in almost all spell checking functions in text editors and command language interfaces. The first spelling correction algorithm based on this technique was proposed by Damerau (1964). Almost at the same time, Levenshtein also developed a similar algorithm. Several other algorithms on edit distance were born thereafter. The edit distance is defined as the minimum number of edit operations required to transform a word to another (Kukich, 1992). These operations are insertion, deletion, substitution and transposition. In most cases, correcting a spelling error requires the insertion, deletion or substitution of a single character, or the transposition of two characters. When a wrong word can be transformed into a dictionary word by inverting one of these operations, the dictionary word is considered a plausible correction. Damerau's algorithm (Damerau, 1964) for edit distance detects spelling errors by comparing words of four to six characters with a list of most frequently used words. When there are multiple candidate words for a given edit distance on a detected word, the first word in the dictionary appearing in alphabetical order is chosen. Levenshtein's algorithm is in the field of dynamic programming and seems to be the most widely used in edit distance computing. Each edit operation is assigned a cost, usually 1 for deletion and insertion and 2 for substitution and transposition. Given a dictionary of  $n$  words, the correction algorithms based on edit distance generally require  $n$  comparisons for each wrong word. To reduce the search time, reversed edit distance technique is used. Another approach used to reduce the number of comparisons involves sorting or partitioning the dictionary according to certain criteria such as alphabetical order, word length, or words occurrences. Many other techniques are also used in spelling errors correction like similarity keys, rules system,  $n$ -grams, probabilistic techniques and neural networks. However, the most widely used technique in errors correction remains edit distance (Hsuan, 2008). It has a time complexity of  $O(nm)$ , with  $n$  and  $m$  the respective sizes of the two compared words. A technique developed by Horst (1993) combining automata and edit

distance was used to quickly find the closest correct word to a wrong word. It has a linear complexity in time relative to the length of the wrong word, regardless of the dictionary size. But the space complexity of the method is exponential.

## Design of a Spell Checker for African Languages

Despite their low level of computerization, African languages have important linguistic resources like bilingual and editorial dictionaries. Although they have a lot of differences, we believe in the possibility of using the little resources they have in common to develop processing tools. Eight West African languages from three different countries are considered in the current study. Five main languages from Niger, two from Mali and one from Burkina Faso are considered. The five languages of Niger are: Fulfulde (ful), Hausa (hau), Kanuri (kau), Songhai-Zarma (son) and Tamashek (tmh). The two languages from Mali are: Bambara (bam) and Soninke (snk) and the one from Burkina Faso is Dyoula (dyu). Though they all use the Latin alphabet for their transcription, each of these languages has its own special characters as shown in Table 1 (Enguehard and Naroua, 2008).

From the literature collected, we believe that the construction of a spell checker will be a step forward towards the computerization of African languages. However, the checker should use minimum resources and efficiently consider the specificities of the concerned languages. Our methodology consists of designing a general tool that uses resources that can easily be obtained from the considered languages despite of their differences with regards to special symbols. Taking into account the linguistic resources available to us, a technique based on a dictionary was found to be more suitable for the design of the spell checker. Although the meaning of a word is contextual, we assume that error detection is independent of the context. An erroneous word is identified by a simple dictionary lookup where the following operations are allowed:

- Add a word to the dictionary
- Check if a word is in the dictionary
- Delete a word from the dictionary

To achieve this, a number of classes are necessary. The class diagram of the entire process is shown in Fig. 1.

The words are represented in form of nodes. Each node has as many links as there are characters in the alphabet. Each valid character string is assigned a value. This may be of any type. It can be used to store information on every word in the dictionary such as definition, grammatical class, translation into another language, etc.

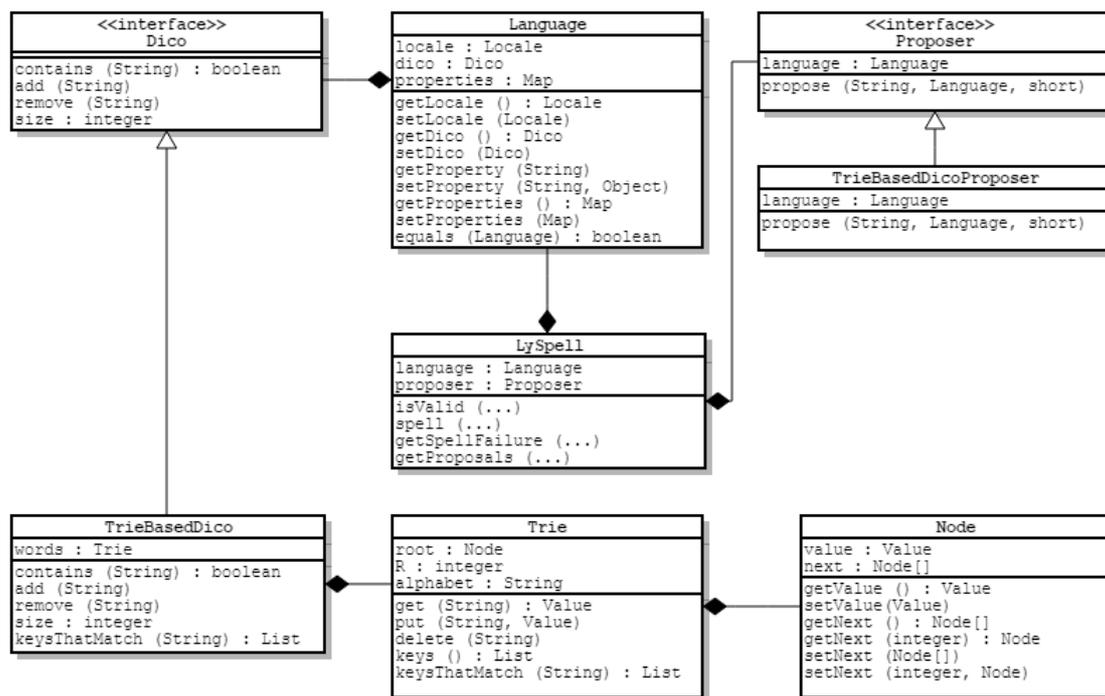


Fig. 1. Global class diagram

Table 1. Special characters used in the alphabets of the studied languages

Name	Sign	Languages	Unicode
Latin letter e with tilde	ẽ	son	U+0065 U+0303
Latin letter i with tilde	ĩ	son	U+0069 U+0303
Latin letter o with tilde	õ	son	U+006F U+0303
Latin letter r short stroke overlay	ṙ	kau	U+0072 U+0335
Latin letter u with tilde	ũ	son	U+0075 U+0303
Latin letter a with tilde	ã	son	U+00E3 or U+61 U+303
Latin letter a with breve	ă	tmh	U+0103 or U+61 U+306
Latin letter ENG	ŋ	bam, ful, son	U+014B
Latin letter s with caron	š	tmh	U+0161 or U+73 U+30C
Latin letter k with hook	ķ	hau	U+0199
Latin letter y with hook	ƿ	ful, hau	U+01B4
Latin letter turned e	ə	kau, tmh	U+01DD or U+259
Latin letter g with caron	ġ	tmh	U+01E7 or U+67 U+30C
Latin letter j with caron	ĵ	tmh	U+01F0 or U+6A U+30C
Latin letter b with hook	ḃ	ful, hau	U+0253
Latin letter open o	ɔ	bam	U+0254
Latin letter d with hook	ḏ	ful, hau	U+0257
Latin letter gamma	γ	tmh	U+0263
Latin letter epsilon	ε	bam	U+025B
Latin letter n with retroflex hook	ɳ	bam, son	U+0272
Latin letter d with dot below	ḏ	tmh	U+1E0D or U+64 U+323
Latin letter l with dot below	ł	tmh	U+1E37 or U+6C U+323
Latin letter s with dot below	ṡ	tmh	U+1E63 or U+73 U+323
Latin letter t with dot below	ṫ	tmh	U+1E6D or U+74 U+323
Latin letter z with dot below	ẓ	tmh	U+1E93 or U+7A U+323

The R attribute of the class Trie is the number of symbols or letters of the alphabet. The characters are represented by indices of *next* array (Node [ ]). Unfortunately, representing characters by indices of *next*

array will set a large value for R. This will inevitably lead to a waste of memory space and additional checks to prevent foreign words from being added to the trie. To avoid this problem, a trick is to find a mapping function

between indices of the *next* array and letters of the alphabet (Robert and Kevin, 2011). That is why the *alphabet* attribute is present in the class *Trie*. It is of type String but it may also be an array of characters. Two additional methods, *toChar* and *toIndex* assure the conversion from indices to characters and vice versa. The *charAt* and the *indexOf* methods of the String class can be effectively used and to make the trick more flexible, we can totally delegate this task to an interface *Alphabet* that defines *toChar* and *toIndex*. The *KeysThatMatch* is another interesting method. Indeed, it allows to search the trie for words that match a given pattern. The patterns used are those with a wildcard, for example a dot ('.'). It is this possibility that we use to implement the reverse editing distance. The *KeysThatMatch* method uses a data structure *List* (a linked list of Strings) to keep the search results. The *List* class has methods to add an item, to verify the existence of an item and to delete an item. Tags may be used to take care of the contextual meanings of a word. It may be very useful for applications like syntax analyzers.

To abstract the implantation of the real dictionary, add flexibility, simplify maintenance and facilitate scalability of the spell checker, an abstract dictionary is represented by a class (*TrieBasedDico*) that implements *Dico* interface (or abstract class). It defines the methods (add, remove, contains) needed to operate on a dictionary. *TrieBasedDico* class is designed by composition from *Trie* class.

The list of candidate words for the correction of an erroneous word is determined in several steps. Once a word is identified as being erroneous, the procedure for determining the type of the error follows. We defined three types of errors (inspired by our research on OpenOffice):

- IS\_NEGATIVE\_WORD: Error caused by the presence of a number or a character not belonging to the alphabet in the word. The word is called negative
- CAPTION\_ERROR: Case Error. This is when a word that should be written with the first letter capitalized is written entirely in lowercase
- SPELLING\_ERROR: represents all other types of spelling errors

The types of errors are short integers encapsulated as static fields in the *LySpell* class. The corrector has two methods for the determination of errors. First, the *getSpellFailure* method which analyzes a given word and returns -1 if the word is correct or one of the three types of errors mentioned above otherwise. Then *isValid* method that checks whether a given word is valid according to the result returned by *getSpellFailure* and spellchecking settings. If *getSpellFailure* returns a value:

- Equal to -1, the word is valid and *isValid* returns true

- Other than -1, the correction parameters are taken into account to determine the validity of the word. For example when you choose not to correct words with numbers and the erroneous word contains digits, *isValid* returns true. This method can be exploited to correct spelling as you type

*currentLanguage* represents the language being supported by the spell corrector. It is an instance of *Language* class. Searching suggestions is performed by *propose* which is an instance of a class that implements the interface *Proposer*.

The method *getProposals* provides correction suggestions for an invalidated word by *isValid* depending on the type of error detected by *getSpellFailure*.

The processed language is represented by the *Language* class. After several attempts, we decided that the dictionary is an attribute of the language and not the reverse. The local attribute of the *Language* class stores information about the processed language. It is of type *Locale* (representation of a language in Java) and provides among others: A two-letter ISO 639-1 code of the language, a two-letter ISO 3166 code of the country as well as the complete names of the language and the country. We use this data for naming resources and for user display. The *properties* attribute is of type Map (mapping key/value) and stores other properties of the language that we use to design the checker and which are not provided by *Locale*. They are currently the alphabet of the language, the special characters in the alphabet, the characters that look like special characters and the punctuation symbols that we divided into two parts: Word separators and end of sentence signs. All the characters of the alphabet are coded in Unicode. The class in charge of finding suggestions implements *Proposer* interface which defines two methods: *isNegativeWord* and *propose*. The *TrieBasedDicoProposer* class uses some features of the alphabet to find candidate words which are found using the reverse edit distance as follows:

- All words having an edit distance equal to 1 with the wrong word are generated by applying edit operations such as insertion, deletion, substitution and transposition. A total of  $60n+28$  words are generated for a wrong word of length  $n$
- Each previously generated word is searched in the trie. If it is there, then it is retained as a possible correction of the erroneous word

The research is conducted by a private method called *proposeByReverseEditDistance*. This method is actually based on *keysThatMatch*. It takes an argument of type *TrieBasedDico* and a word or a pattern and returns the result as an array of Strings. Methods that perform

editing operations on a given word are provided by the *StringTools* class which consists of tools shared by different classes. The minimum edit distance is used to rank the suggested words. Those who are closest to the wrong word are placed at the top of the list. To implement that, a comparator was designed.

## Implementation

In this article, a software tool has been proposed for processing eight West African languages from three different countries. It has been designed as a spell checker for which only a dictionary and the alphabet of a language are needed as linguistic resources. A prototype of the designed software tool has been implemented in Java under Linux environment. For the sake of application, a dictionary (Mijinguini, 2003) and the official alphabet of the Hausa language of Niger are used. The developed tool was tested as a standalone program through a text editor designed for that purpose and as an extension for the OpenOffice suite. Figure 2 and 3 respectively show the dialog box for spell checking and error correction in OpenOffice writer. However, in OpenOffice, the Hausa options are only Nigeria and Ghana and we are limited to select one of them. From the results, we observe that:

- It is possible, from minimum language resources and proven techniques to develop automatic processing tools for African languages
- Only a dictionary as language resource and the alphabet of an African language are needed to develop a spell checker
- The specificities of the African languages can be efficiently handled
- Although they contain special characters, the dictionaries of African languages can efficiently be represented by tries

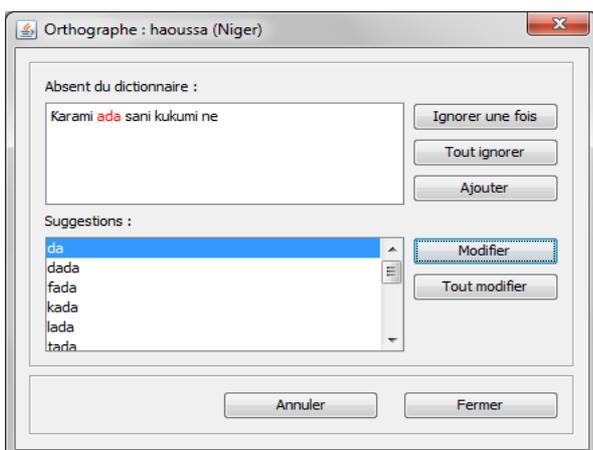


Fig. 2. Dialog box for error correction

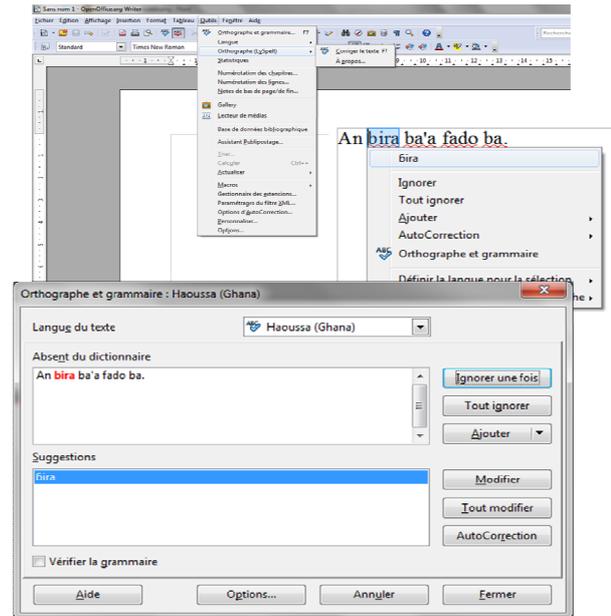


Fig. 3. Error correction

## Conclusion

This work is a contribution to the automatic processing of African languages. Although it may be necessary to improve the performance of the designed tool, we believe that the results we have obtained will add value to the computerization of African languages and contribute to their effective use in institutions of education and on media.

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## Author's Contributions

The authors have equally contributed to the research work and writing of the article.

## Ethics

This article is original and contains unpublished material. The corresponding author confirms that all authors have read and approved the manuscript and no ethical issues are expected to arise after its publication.

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