# Framework of Regression-Based Graph Matrix Analysis Approach in Multi-Relational Social Network Problem 

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#### Abstract

Community mining is one of the major directions in social network analysis. Social network analysis has attracted much attention in recent years. Most of the existing methods on community mining assume that there is only one kind of relation in the network and moreover, the mining results are independent of the users' needs or preferences. However, in reality, there exist multiple, heterogeneous social networks, each representing a particular kind of relationship and each kind of relationship may play a distinct role in a particular task. Thus mining networks by assuming only one kind of relation may miss a lot of valuable hidden community information and may not be adaptable to the diverse information needs from different users. In this research, we systematically analyze the problem of mining hidden communities on heterogeneous social networks. Based on the observation that different relations have different importance with respect to a certain query, we propose a method for learning an optimal linear combination of these relations which can best meet the user's expectation. With the obtained relation, better performance can be achieved for community mining.


Key words: Community mining, social network analysis, hidden community information

## INTRODUCTION

Social Network Analysis (SNA) is the mapping and measuring of relationships and flows between people, groups, organizations, computers, or other information/knowledge processing objects. Social network analysis as a theme has been studied for years. The classic paper of Milgram ${ }^{[23]}$ might be one of the first works on SNA. It estimates that every person in the world is only six edges away from every other. It sets the stage for investigations into social networks and algorithmic aspects of social networks. Many recent efforts try to leverage social networks for diverse purposes, such as expertise location ${ }^{[19,19]}$, mining the network value of customers ${ }^{[11]}$ and discovering shared interests ${ }^{[26]}$.

Previous work in sociology and statistics has suffered from the lack of data and focused on very small networks, typically in the tens of individuals ${ }^{[29]}$. With the web growing, much potential social network data are available and a lot research efforts have been put on dealing with such data.

Schwartz and Wood mined social relationships from email logs ${ }^{[31]}$. The ReferralWeb project ${ }^{[19]}$ is proposed to mine a social network from a wide variety of web data and use it to help individuals find experts who could answer their questions. Adamic and Adar
tried to discover the social interactions between people from the information on their homepages ${ }^{[1]}$.

Agrawal et al. analyzed the social behavior of the people on the newsgroups ${ }^{[2]}$. Moreover, the web itself can be acually viewed as a large social network.
The well-known link analysis algorithms, such as Google's PageRank ${ }^{[24,16,5,6]}$ and Kleigberg's HITS algorithm ${ }^{[20]}$, can be seen as social network analysis on the web.

Community mining: With the growth of the web, community mining has attracted increasing attention. A lot of work has been done at mining the implicit communities of web pages ${ }^{[14,22,7,12,27,13]}$, scientific literature from the Web ${ }^{[31]}$ and document citation database ${ }^{[25]}$.

In principle, a community can be simply defined as a group of objects sharing some common properties. Community mining has many similar properties to the graph-cut problem. Kumar et al. used the bipartite graph concept to find the core of the community and then expanded the core to get the desired community ${ }^{[22]}$. Flake et al. applied the maximum-flow and minimumcut framework on the community mining ${ }^{[12]}$. The authority-and-hub idea ${ }^{[20]}$ was also used in the community mining ${ }^{[14,21,8]}$ and has several extensions ${ }^{[9]}$. The idea of frequent itemset in association rule mining has also been used in community mining ${ }^{[37]}$.

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Generally speaking, both social network analysis and community mining can be seen as graph mining. The community mining can be thought of as sub-graph identification. Previous work on graph mining can be found in ${ }^{[10,28,30]}$. Almost all the previous techniques on graph mining and community mining are based on a homogenous graph, i.e., there is only one kind of relationship between the objects. However, in real social networks, there are always various kinds of relationships between the objects. To deal with this problem, we focus in this research on multi-relational community mining.

We begin on next section about relation extraction and follow with discussion about multi-relational social network and close with conclusion.

## RELATION EXTRACTION

Here we begin with a detailed analysis of the relation extraction problem followed by two algorithms for two cases.

The problem: A typical social network likely contains multiple relations. Different relations can be modelled by different graphs. These different graphs reflect the relationship of the objects from different views.

For the problems of community mining, these different relation graphs can provide us with different communities.

In multi-relational social network, community mining should be dependent on the user's example (or information need). A user's query can be very flexible. Since previous community mining techniques only focus on single relational network and are independent of the user's query, they cannot cope with such a complex situation.

In this research, we focus on the relation extraction problem in multi-relational social network.

The community mining based on the extracted relation graph is more likely to meet the user's information need. For relation extraction, it can be either linear or nonlinear. Due to the consideration that in real world applications it is almost impossible for a user to provide sufficient information, nonlinear techniques tend to be unstable and may cause overfitting problems. Therefore, here we only focus on linear techniques.

This problem of relation extraction can be mathematically defined as follows. Given a set of objects and a set of relations which can be represented by a set of graphs $G_{i}\left(V, E_{i}\right), i=1, \ldots, n$, where $n$ is the number of relations, V is the set of nodes (objects) and $\mathrm{E}_{\mathrm{i}}$ is the set of edge with respect to the i-th relation. The
weights on the edges can be naturally defined according to the relation strength of two objects. We use Mi to denote the weight matrix associated with $G_{i}, i=1, \ldots, n$. Suppose there exists a hidden relation represented by a graph $\hat{G}(V, \hat{E})$, and $\hat{M}$ denotes the weight matrix associated with $\hat{G}$. Given a set of labeled objects $X=$ $\left[x_{1}, \ldots, x_{m}\right]$ and $y=\left[y_{1}, \ldots, y_{m}\right]$ where $y_{j}$ is the label of $x_{j}$ (Such labeled objects indicate partial information of the hidden relation $\hat{G}$ ), find a linear combination of the weight matrices which can give the best estimation of the hidden matrix $\hat{M}$.

A regression-based algorithm: The basic idea of our algorithm is trying to find an combined relation which makes the relationship between the intra-community examples as tight as possible and at the same time the relationship between the inter-community examples as loose as possible.

For each relation, we can normalize it to make the biggest strength (weight on the edge) be 1 .

Thus we construct the target relation between the labeled objects as follows:

$$
\tilde{\mathrm{M}}_{\mathrm{ij}}= \begin{cases}1, \text { example } \mathrm{i} \text { and example } \mathrm{j} \text { have the same label } \\ 0, & \text { otherwise }\end{cases}
$$

where, $\tilde{M}$ is a $m \times m$ matrix and $\tilde{\mathrm{M}}_{\mathrm{ij}}$ indicates the relationship between examples i and j . Once the target relation matrix is built, we aim at finding a linear combination of the existing relations to optimally approximate the target relation in the sense of $\mathrm{L}_{2}$ norm. Sometimes, a user is uncertain if two objects belong to the same community and can only provide the possibility that two objects belong to the same community. In such case, we can define as follows $\tilde{M}$.
$\tilde{\mathrm{M}}_{\mathrm{ij}}=\operatorname{Prob}\left(\mathrm{X}_{\mathrm{i}}\right.$ and $\mathrm{X}_{\mathrm{j}}$ belong to the same community)
Let $\mathrm{a}=\left[\mathrm{a}_{1}, \mathrm{a}_{2}, \ldots, \mathrm{a}_{n}\right]^{\mathrm{T}} \in \mathrm{R}^{\mathrm{n}}$ denote the combination coefficients for different relations. The approximation problem can be characterized by solving the following optimization problem:

$$
\begin{equation*}
\mathrm{a}^{\text {opt }}=\arg \min _{\mathrm{a}}\left\|\tilde{\mathrm{M}}-\sum_{\mathrm{i}=1}^{\mathrm{n}} \mathrm{a}_{\mathrm{i}} \mathrm{M}_{\mathrm{i}}\right\|^{2} \tag{1}
\end{equation*}
$$

This can be written as a vector form. Since the matrix $M_{m \times m}$ is symmetric, we can use a $m(m-1) / 2$ dimensional vector v to represent $i$. The problem (1) is equivalent to:

$$
\begin{equation*}
a^{\text {opt }}=\arg \min _{a}\left\|\tilde{\mathrm{~V}}-\sum_{\mathrm{i}=1}^{\mathrm{n}} \mathrm{a}_{\mathrm{i}} \mathrm{~V}_{\mathrm{i}}\right\|^{2} \tag{2}
\end{equation*}
$$

Equation 2 is actually a linear regression problem ${ }^{[15]}$. From this point of view, the relation extraction problem is interpreted as a prediction problem. Once the combination coefficients are computed, the hidden relation strength between any object pair can be predicted.

In real applications, the user does not need to specify the relationships between any pair of objects. That is, the vector $v$ need not to be $m(m-1) / 2$ dimensional. We assume that v is a k-dimensional vector in the following.

Let us first consider the simplest case that:

$$
\begin{equation*}
\sum_{i=1}^{n} a_{i} v_{i}=\tilde{v} \tag{3}
\end{equation*}
$$

We define:

$$
\begin{equation*}
\mathrm{V}=\left[\mathrm{v}_{1}, \mathrm{v}_{2}, \ldots, \mathrm{v}_{\mathrm{n}}\right] \tag{4}
\end{equation*}
$$

Thus, Eq. 3 can be rewritten as follows:

$$
\mathrm{Va}=\tilde{\mathrm{v}}
$$

Suppose the rank of V is $\min (\mathrm{k}, \mathrm{n})$. We have the following facts:

- When $\mathrm{k}<\mathrm{n}$, the set of solutions to Eq. 4 forms a ( $\mathrm{n}-\mathrm{k}$ ) dimensional vector subspace
- When $\mathrm{k}=\mathrm{n}$, there is a unique solution to Eq. 4
- When $\mathrm{k}>\mathrm{n}$, there is no solution to Eq. 4

In the first two cases, we get a solution with perfect match (The minimization error is zero). Note that, the value of $k$ reflects the quantity of the user's information needs. k is small when the query submitted by the user is simple.

With a complex query, $k$ can be larger than $n$. In this case, the optimal solution to (2) is obtained when the derivative of this objective function with respect to a is zero, i.e.
$\frac{\partial\left\|\tilde{v}-\sum_{i=1}^{n} a_{i} v_{i}\right\|^{2}}{\partial a_{i}}=0$ for $i=1,2, \ldots, n$

By some algebraic steps, we have:
$\frac{\partial\left\|\tilde{v}-\sum_{i=1}^{n} a_{i} v_{i}\right\|^{2}}{\partial a_{i}}=0$
$\frac{\partial\left[\left(\tilde{v}-\sum_{i=1}^{n} a_{i} v_{i}\right)^{T}\left(\tilde{v}-\sum_{i=1}^{n} a_{i} v_{i}\right)\right]}{\partial a_{i}}=0$
$\frac{\partial\left(\tilde{v}-\sum_{i=1}^{n} a_{i} v_{i}\right)^{T}}{\partial a_{i}}\left(\tilde{v}-\sum_{i=1}^{n} a_{i} v_{i}\right)+\left(\tilde{v}-\sum_{i=1}^{n} a_{i} v_{i}\right)^{T} \frac{\partial\left(\tilde{v}-\sum_{i=1}^{n} a_{i} v_{i}\right)}{\partial a_{i}}=0$
$2\left\{\frac{\partial\left(\tilde{v}-\sum_{i=1}^{n} a_{i} v_{i}\right.}{\partial a_{i}}\right\}^{2}\left(\tilde{v}-\sum_{i=1}^{n} a_{i} v_{i}\right)=0$
$v_{i}^{T}\left(\tilde{v}-\sum_{i=1}^{n} a_{i} v_{i}\right)=0$ for $i=1,2, \ldots, n$.
$v^{T}\left(\tilde{v}-\sum_{i=1}^{n} a_{i} v_{i}\right)=0$
$v^{T}(\tilde{v}-V a)=0$
$v^{T} \mathrm{Va}=\mathrm{V}^{\mathrm{T}} \tilde{v}^{\sim}$

Since the matrix V has full rank as we assumed, i.e., $\operatorname{rank}(\mathrm{V})=\min (\mathrm{k}, \mathrm{n})$, the matrix $\mathrm{V}^{\mathrm{T}} \mathrm{V}$ is invertible and the optimal solution to (2) is $a^{\text {opt }}=\left(V^{T} V\right)^{-1} V^{T} \tilde{v}$.

When the matrix V is rank deficiency, i.e., $\operatorname{rank}(\mathrm{V})$ $<\min (k, n)$, there will be multiple solutions with the same minimization value. In such case, we can choose the a with minimum norm as our solution ${ }^{[4]}$.

The objective function (2) models the relation extraction problem as an unconstrained linear regression problem. One of the advantages of the unconstrained linear regression is that, it has a close form solution and is easy to compute. However, researches on linear regression problem show that in many cases, such unconstrained least squares solution might not be a satisfactory solution and the coefficient shrinkage technique should be applied based on the following two reasons ${ }^{[15]}$.

Prediction accuracy: The least-squares estimates often have low bias but large variance ${ }^{[15]}$.

The overall relationship prediction accuracy can sometimes be improved by shrinking or setting some coefficients to zero. By doing so we sacrifice a little bit of bias to reduce the variance of the predicted relation strength and hence may improve the overall relationship prediction accuracy.

Interpretation: With a large number of explicit (base) relation matrices and corresponding coefficients, we often would like to determine a smaller subset that exhibit the strongest effects.

In order to get the big picture, we are willing to sacrifice some of the small details.

Such consideration can be shown in the following example. Suppose we have a user query $\left(\mathrm{o}_{1}, \mathrm{o}_{2}, \mathrm{o}_{3}, \mathrm{o}_{4}\right.$, $\mathrm{o}_{5}$ ), where $\mathrm{o}_{1}, \mathrm{o}_{2}$ and $\mathrm{o}_{3}$ belong to one community, but $\mathrm{o}_{4}$ and $\mathrm{o}_{5}$ belong to another.

The target relation network can be constructed as:

|  | $\mathrm{O}_{1}$ | $\mathrm{O}_{2}$ | $\mathrm{O}_{3}$ | $\mathrm{O}_{4}$ | $\mathrm{O}_{5}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{O}_{1}$ | $*$ | 1 | 1 | 0 | 0 |
| $\mathrm{O}_{2}$ | 1 | $*$ | 1 | 0 | 0 |
| $\mathrm{O}_{3}$ | 1 | 1 | $*$ | 0 | 0 |
| $\mathrm{O}_{4}$ | 0 | 0 | 0 | $*$ | 1 |
| $\mathrm{O}_{5}$ | 0 | 0 | 0 | 1 | $*$ |

The $*$ in the relation matrix means that we do not consider the self-relation strength. The four basic relation matrices (corresponding to these objects) in the social networks are shown in Fig. 1.

We can find that $0 M_{1}+10 M_{2}+10 M_{3}+10 M_{4}$ can exactly match the example relation matrix. However, such extracted relation might not a good approximation to the hidden relation on the whole object set. The relation M1 is more likely to be what the user desires. This is exactly the problem of unconstrained linear regression. We need to use some coefficient shrinkage techniques to solve such problem ${ }^{[17]}$.

Thus, for each relation network, we normalize all the weights on the edges in the range $[0,1]$.

And, we put a constraint: $\sum_{i=1}^{n} a_{i}^{2} \leq 1$ on the objective function (2). Finally, our algorithm tries to solve the following minimization problem,

$$
\begin{equation*}
\mathrm{a}^{\mathrm{opt}}=\arg \min _{\mathrm{a}}\left\|\tilde{\mathrm{v}}-\sum_{\mathrm{i}=1}^{\mathrm{n}} \mathrm{a}_{\mathrm{i}} \mathrm{e}_{\mathrm{i}}\right\| \tag{5}
\end{equation*}
$$

Subject to $\sum_{i=1}^{n} \mathrm{a}_{\mathrm{i}}^{2} \leq 1$
Such a constrained regression is called Ridge Regression ${ }^{[15]}$ and can be solved by some numerical methods ${ }^{[4]}$. When we use such constrained relation extraction, the coefficients of the extracted relation for the above example are $1,0,0,0$. This shows that our constrained relation extraction can really solve the problem.

A MinCut-based algorithm: In the last subsection, we have presented a general method for exacting the

|  | $\mathrm{O}_{1}$ | $\mathrm{O}_{2}$ | $\mathrm{O}_{3}$ | $\mathrm{O}_{4}$ | $\mathrm{O}_{5}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{O}_{1}$ | $*$ | 0.8 | 0.7 | 0 | 0 |
| $\mathrm{O}_{2}$ | 0.8 | $*$ | 0.9 | 0 | 0 |
| $\mathrm{O}_{3}$ | 0.7 | 0.9 | $*$ | 0 | 0 |
| $\mathrm{O}_{4}$ | 0 | 0 | 0 | $*$ | 0.6 |
| $\mathrm{O}_{5}$ | 0 | 0 | 0 | 0.6 | $*$ |

(a) Relation $\mathrm{M}_{1}$

|  | $\mathrm{O}_{1}$ | $\mathrm{O}_{2}$ | $\mathrm{O}_{3}$ | $\mathrm{O}_{4}$ | $\mathrm{O}_{5}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{O}_{1}$ | $*$ | 0 | 0.1 | 0 | 0 |
| $\mathrm{O}_{2}$ | 0 | $*$ | 0 | 0 | 0 |
| $\mathrm{O}_{3}$ | 0.1 | 0 | $*$ | 0 | 0 |
| $\mathrm{O}_{4}$ | 0 | 0 | 0 | $*$ | 0 |
| $\mathrm{O}_{5}$ | 0 | 0 | 0 | 0 | $*$ |

(b) Relation $\mathrm{M}_{2}$

|  | $\mathrm{O}_{1}$ | $\mathrm{O}_{2}$ | $\mathrm{O}_{3}$ | $\mathrm{O}_{4}$ | $\mathrm{O}_{5}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{O}_{1}$ | $*$ | 0.1 | 0 | 0 | 0 |
| $\mathrm{O}_{2}$ | 0.1 | $*$ | 0 | 0 | 0 |
| $\mathrm{O}_{3}$ | 0.1 | 0 | $*$ | 0 | 0 |
| $\mathrm{O}_{4}$ | 0 | 0 | 0 | $*$ | 0.1 |
| $\mathrm{O}_{5}$ | 0 | 0 | 0 | 0.1 | $*$ |

(c) Relation $\mathrm{M}_{3}$

|  | $\mathrm{O}_{1}$ | $\mathrm{O}_{2}$ | $\mathrm{O}_{3}$ | $\mathrm{O}_{4}$ | $\mathrm{O}_{5}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{O}_{1}$ | $*$ | 0 | 0 | 0 | 0 |
| $\mathrm{O}_{2}$ | 0 | $*$ | 0.1 | 0 | 0 |
| $\mathrm{O}_{3}$ | 0 | 0.1 | $*$ | 0 | 0 |
| $\mathrm{O}_{4}$ | 0 | 0 | 0 | $*$ | 0.1 |
| $\mathrm{O}_{5}$ | 0 | 0 | 0 | 0.1 | $*$ |

(d) Relation $\mathrm{M}_{4}$

Fig. 1: The four basic relation matrices corresponding to the examples
hidden relation based on regression model. However, this method may fail when the examples provided by the user belong to only one community, which is referred to single community issue in the rest of this research. We provide an intuitive example in the following.

Suppose we have a user query $\left(\mathrm{o}_{1}, \mathrm{o}_{2}, \mathrm{o}_{3}, \mathrm{o}_{4}, \mathrm{o}_{5}\right)$, which belong to the same community. In the following two relations shown in Fig. 2a and b, regression model would prefer the relation $M_{1}$, since the higher connectivity between $\mathrm{o}_{1}, \mathrm{o}_{2}, \mathrm{o}_{3}, \mathrm{o}_{4}$ achieves a lower square error to the target relation. However, in relation $\mathrm{M}_{1}$, the connectivity between $\mathrm{o}_{5}$ and the other four examples are very weak. As can be seen, the connectivity in $\mathrm{M}_{2}$ is much more uniform than that in

|  | $\mathrm{O}_{1}$ | $\mathrm{O}_{2}$ | $\mathrm{O}_{3}$ | $\mathrm{O}_{4}$ | $\mathrm{O}_{5}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{O}_{1}$ | $*$ | 0 | 0 | 0 | 0 |
| $\mathrm{O}_{2}$ | 0 | $*$ | 0.1 | 0 | 0 |
| $\mathrm{O}_{3}$ | 0 | 0.1 | $*$ | 0 | 0 |
| $\mathrm{O}_{4}$ | 0 | 0 | 0 | $*$ | 0.1 |
| $\mathrm{O}_{5}$ | 0 | 0 | 0 | 0.1 | $*$ |

(a) Relation $\mathrm{M}_{1}$

|  | $\mathrm{O}_{1}$ | $\mathrm{O}_{2}$ | $\mathrm{O}_{3}$ | $\mathrm{O}_{4}$ | $\mathrm{O}_{5}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{O}_{1}$ | $*$ | 0.4 | 0.4 | 0.5 | 0.5 |
| $\mathrm{O}_{2}$ | 0.4 | $*$ | 0.3 | 0.2 | 0.4 |
| $\mathrm{O}_{3}$ | 0.4 | 0.3 | $*$ | 0.3 | 0.3 |
| $\mathrm{O}_{4}$ | 0.5 | 0.2 | 0.3 | $*$ | 0.4 |
| $\mathrm{O}_{5}$ | 0.5 | 0.4 | 0.3 | 0.4 | $*$ |

(b) Relation $\mathrm{M}_{2}$

Fig. 2: Two existing relations
$\mathrm{M}_{1}$ while it has comparable strength. Therefore, $\mathrm{M}_{2}$ should be a better choice for this user query. Unfortunately, the square error of $\mathbf{M}_{2}$ is larger than that of $\mathrm{M}_{1}$. This shows that the regression model may fail in such a case.

In order to deal with the single community issue, we need to take into account the weakest connection in the extracted relation. By graph theory, the value of the minimum cut on the graph can be used to evaluate the tightness of the graph.

Let $G$ denote a weighted graph with weight matrix M . Let m denote the number of vertices.

A cut on the graph is defined as a set of edges which separates the vertices into two disconnected groups denoted by $A$ and $B$ such that $A \cap B=\varnothing$ and $A \cup B=D$. Thus, the value of the cut is: $\operatorname{cut}(\mathrm{G})=\sum_{\mathrm{i} \in \mathrm{A}} \sum_{\mathrm{j} \in \mathrm{B}} \mathrm{M}(\mathrm{i}, \mathrm{j})$

It is easy to see that there are totally $2^{\mathrm{m}}-2$ different cuts. Let $\operatorname{cutk}(G)=\left(A_{k}, B_{k}\right)$ denote the k-th cut. The minimum cut is defined as: $\min \operatorname{cut}(G)=\min _{\mathrm{k}}\left\{\operatorname{cut}_{\mathrm{k}}(\mathrm{G})\right\}$

If a graph can be easily cut into two subgraphs, it has a small minimum cut value. As an extreme case, the minimum cut value of a disconnected graph is 0 . Naturally, the optimal extracted relation graph should have a large minimum cut value. Thus, for single community issue, we try to extract the optimal relation graph by maximizing its minimum cut value.

Let $G_{i}, i=1, \ldots, n$, denote the existing relation graphs defined only on the user query examples and Mi denote the corresponding weight matrices. Let $\mathrm{a}=\left[\mathrm{a}_{1}\right.$, $\left.a_{2}, \ldots, a_{n}\right]^{T} \in R_{n}$ denote the combination coefficients for different graphs. Thus $M=\sum_{i=1}^{n} a_{i} M_{i}$ is the weight matrix
of the combined relation graph $G$. Let mincut( $G$ ) denote the minimum cut value of G. Our objective function can be written as follows:

$$
\begin{equation*}
a^{\text {opt }}=\arg \min _{a}\left\|\tilde{M}-\sum_{i=1}^{n} a_{i} M_{i}\right\|^{2} \tag{6}
\end{equation*}
$$

Generally, the minimum cut problem is an NP-hard problem. Thus the optimization problem (6) cannot be easily solved. However, in our problems, the number of examples provided by the user is usually small. That is, $m$ is small, typically less than 10 . Thus we can use linear programming techniques to solve the optimization problem (6) by the following derivation: mincut(G)

$$
\begin{aligned}
& =\min _{1 \leq k \leq 2^{m}-2}\left\{\operatorname{cut}_{\mathrm{k}}(\mathrm{G})\right\} \\
& =\min _{1 \leq \mathrm{k} \leq 2^{m}-2}\left\{\sum_{i \in \mathrm{~A}(\mathrm{k})} \sum_{\mathrm{j} \in \mathrm{~B}(\mathrm{k})}\left(\sum_{\mathrm{l}=1} \mathrm{a}_{1} \mathrm{M}_{1}(\mathrm{i}, \mathrm{j})\right)\right\} \\
& =\min _{1 \leq \mathrm{k} \leq 2^{\mathrm{m}}-2}\left\{\sum_{i \in \mathrm{~A}(\mathrm{k})} \mathrm{a}_{1}\left(\sum_{\mathrm{i} \in \mathrm{~A}(\mathrm{k})} \sum_{j \in \mathrm{~B}(\mathrm{k})} \mathrm{M}_{1}(\mathrm{i}, \mathrm{j})\right)\right\} \\
& =\min _{1 \leq \mathrm{k} \leq 2^{\mathrm{m}}-2}\left\{\sum_{i \in \mathrm{~A}(\mathrm{k})} \mathrm{a}_{1} \cdot \operatorname{cut}_{\mathrm{k}}\left(\mathrm{G}_{1}\right)\right\}
\end{aligned}
$$

Let $\mathrm{v}=\operatorname{mincut}(\mathrm{G})$. The optimization problem in Eq. 6 can be reduced to the following linear programming problem:
Max v
$\sum_{\mathrm{l}=1}^{\mathrm{n}} \mathrm{a}_{1} \cdot \operatorname{cut}_{\mathrm{k}}\left(\mathrm{G}_{1}\right)-\mathrm{v} \geq 0, \quad\left(1 \leq \mathrm{k} \leq 2^{\mathrm{m}}-2\right)(*)$
$\sum_{\mathrm{l}=1}^{\mathrm{n}} \mathrm{a}_{1}=1$
$\mathrm{a}_{1} \geq 0, \quad(1 \leq 1 \leq \mathrm{n})$

With the constraints $\left(^{*}\right)$, v is guaranteed to be the minimum cut value and by maximizing v we can obtain the optimal combination coefficients ai. The number of constraints in this problem is $2^{m}-2+n+1$, where $m$ is the number of user-provided examples which is usually less than 10 and n is the number of existing relations. The above problem can be efficiently solved by linear programming techniques ${ }^{[3]}$.

The proposed regression based algorithm and the MinCut based algorithm are used under different situations. When a user provides multiple community examples, regression-based algorithm can be used to find the best combination; when he provides single
community examples, MinCut-based algorithm can be used.

## DISCUSSION

Since mining hidden communities in heterogeneous networks represents a promising research direction, there are many issues that need to be discussed. Here we focus on the problem solving philosophy.

First, one may wonder the complexity at comprehension and combination of multiple social networks in the analysis. We do agree that multiple social networks form complex, multiple, interrelated graphs and with the massive amount of data mounting, it is challenging for anyone to grasp the whole picture of such dynamic, evolving social networks and work out a balanced combination of multiple networks for a particular user query. However, such multiple networks do exist and it is inappropriate to blindly merge them into one since different networks plays different roles in particular queries, as shown in our experiments. Therefore, we believe that developing new multinetwork mining algorithms to dynamically combine multiple relevant networks to form combined virtual networks based on user's example queries is a new and appropriate problem solving methodology.

Second, since it is difficult for a user to comprehend the whole picture of numerous social networks, one may wonder how a user is able to pose high-quality queries. Based on our experience, although it is difficult for a user to comprehend the overall multiple networks, a user usually has good knowledge on a small set of examples (such as influential researchers, movie/sport stars, big companies, or popular commodities). Such firm grasp of a small set of examples is often sufficient to pose intelligent queries, learn additional facts and form informative combined networks.

Third, one may wonder how to comprehend the answers returned from such a network analysis.

Since a derived hidden network is a weighted matrix as a combination of multiple existing networks, it is often difficult to understand the minor weight differences in the results. However, the real essence is at the new facts derived from such hidden networks and their associated rankings. This resembles Google-like keyword-based Web search. It is not so crucial to understand the derived Web linkage weighting and claim it is optimal. However, the return of quality rankings on the interesting results demonstrate its utility.

## CONCLUSION

Different from most social network analysis studies, we assume that there exist multiple, heterogeneous social networks and the sophisticated combinations of such heterogeneous social networks may generate important new relationships that may better fit user's information need. Therefore, our approach to social network analysis and community mining represents a major shift in methodology from the traditional one, a shift from single-network, userindependent analysis to multi-network, user-dependant and query-based analysis. Our argument for such a shift is clear: multiple, heterogeneous social networks are ubiquitous in the real world and they usually jointly affect people's social activities.

Based on such a philosophy, we worked out a new methodology for relation extraction and proposed two algorithms in different situations. With such querydependent relation extraction and community mining, fine and subtle semantics are captured effectively. Our discussion also shows it is expected that the querybased relation extraction and community mining would give rise to a lot of potential new applications in social network analysis.

There are a lot of issues that need to be studied further. First, our approach adopts a regression-based graph matrix analysis approach. There are potentially many other approaches that can be explored and compared with this approach. We will expect that future studies may propose even more powerful approaches in relation extraction than what is proposed here.

Second, our relation extraction algorithm has made a lot of simplifications in the analysis. In general, links within the same network and among different networks may carry different weights. For example, one can imagine that the links among co-author networks should be inherently stronger than those among coproceedings since average size (\# of links) in the coauthor group is much smaller than that in the coproceedings group. This is not considered in our simple model. Thus we expect the prediction power will be substantially enhanced if such information is incorporated in the new algorithm.

Third, our query model considers only one simple group of nodes (such as researchers). A more powerful query model may involve and, or, not operators on those groups. For example, one may like to find those who co-attend the same conference but never coauthored a paper using the not operator. This will be useful for finding referees for conference submissions. These issues may form an exciting frontier for future research.

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